

Random Numbers

Random number generator (Uniform)

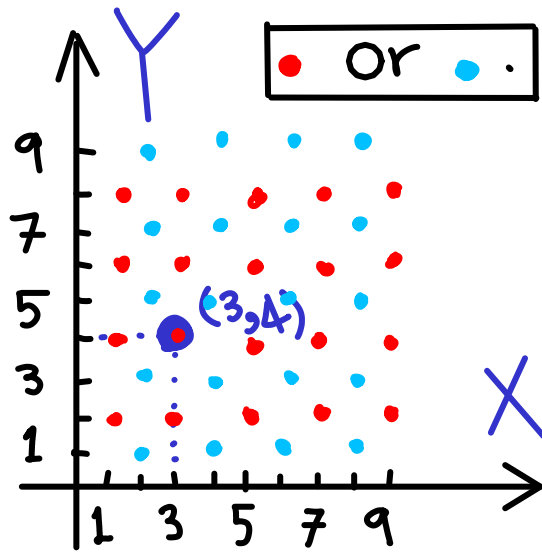
$$X_{n+1} = \left(\underset{\substack{\downarrow \\ \text{odd}}}{a} X_n + \underset{\substack{\downarrow \\ \text{odd}}}{c} \right) \bmod \underset{\substack{\downarrow \\ 2^{32}}}{m} \rightarrow m \text{ should be a big number}$$

Suppose that we want to have a random generator to produce random numbers between 0 and 9. How?

If the random number generated by the above relation is: 581
one can take the right digit of this number. Is this a good way?

answer:

No. Because if X_n is an even number then X_{n+1} will be an odd number, and vice versa. Therefore, if we take the right digit of X_{n+1} then once we have an even number and right after that we will have an odd number. Consequently, there is a correlation between the produced numbers in $[0, 9]$.



If we take two numbers x and y , and make a vector (x, y) . Then by plotting these vectors we can not cover all space in 2D between $[0, 9]$. For example $(3, 4)$ has been shown.

One way is to take the digit before the most right digit, i.e. 5~~8~~1.

The most left digit is not good as well. Because its distribution is not uniform and is logarithmic.

Suffling algorithm (suitable for very long simulations)

$$X_{n+1} = (aX_n + c) \bmod m \quad \rightarrow$$

$$\begin{bmatrix} x_0 \\ x_1 \\ \vdots \\ x_{498} \\ x_{499} \end{bmatrix}$$

1. For example we take $m' = 500$, and we use also the following random number generator.

$$Y_{n+1} = (a'Y_n + c') \bmod m'. \text{ Then } 0 \leq Y_{n+1} < 500.$$

2. We produce m' random numbers by using $X_{n+1} = \dots$ and put them in an array.

3. We produce the random number Y_{n+1} , and then take Y_{n+1} th element of the array as a desired random number.

4. We make the array empty.

5. Go to 2.

```

#include <stdio.h>
#include <stdlib.h>
#include <time.h>

int main () {
    int i, n, rr, seed_number;
    time_t t;
    //
    seed_number = time(&t);
    //seed_number = 123456789;
    n = 5;

    /* Intializes random number generator */
    //srand((unsigned) time(&t));
    srand((unsigned) seed_number);
    printf("%ld\n\n\n\n", time(&t));

    /* Print 5 random numbers from 0 to 49 */
    for( i = 0 ; i < n ; i++ ) {
        rr = rand();
        printf("%d\n", rr % 100);
        printf("%d\n", rr );
    }

    return(0);
}

```

```

2115143165
65
665142609
9
462857371
71
813125836
36
632321020
20

```

```

#include <stdio.h>
#include <stdlib.h>
#include <time.h>

int main () {
    int i, n, rr, seed_number;
    time_t t;
    //
    n = 5;

    /* Intializes random number generator */
    //printf("%ld\n\n\n\n", time(&t));

    /* Print 5 random numbers from 0 to 99 */
    for( i = 0 ; i < n ; i++ ) {
        seed_number = time(&t);
        srand((unsigned) seed_number);

        rr = rand();
        printf("%d\n", rr );
        printf("%d\n", rr % 100);
    }

    return(0);
}

```

```

2065990336
36
2065990336
36
2065990336
36
2065990336
36
2065990336
36

```

- Many phenomena in physics reveal Gaussian random distribution. For example velocity distribution of a gas is Gaussian.
- Therefore, we need to know how non-uniform random numbers can be generated.

Central Limit Theorem (in stat. mech.)

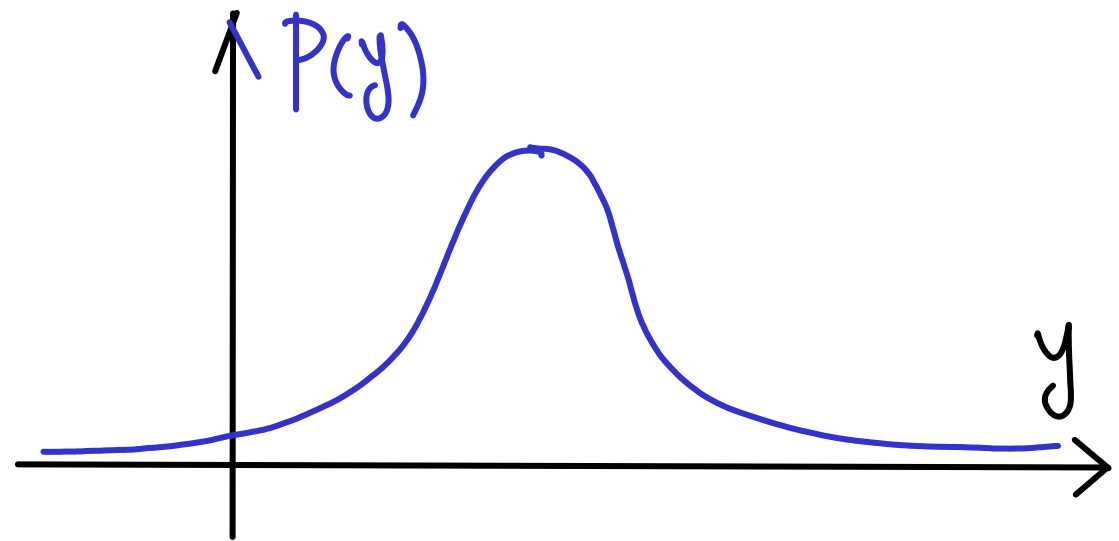
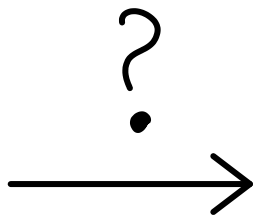
$$X_{n+1} = (a X_n + c) \text{ mod } m$$

If $\begin{cases} X_i \text{ s are uncorrelated and random.} \\ y = \frac{1}{N} \sum_{i=1}^N X_i \text{ (If } N \text{ is big enough)} \end{cases} \Rightarrow$

$$P(y) = \frac{1}{\sqrt{2\pi} \tilde{\sigma}_y} e^{-\frac{(y - \langle y \rangle)^2}{2 \tilde{\sigma}_y^2}}$$

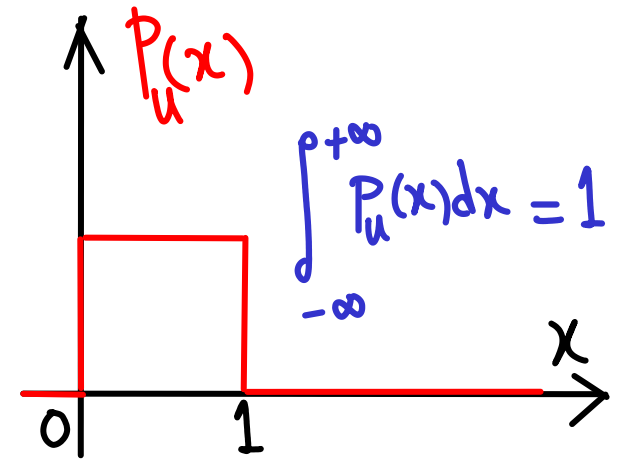
$$\langle y \rangle = \langle x \rangle \quad \tilde{\sigma}_y = \tilde{\sigma}_x / \sqrt{N}$$

How a non-uniform probability distribution can be produced from a uniform probability distribution?



$$P(x) = P_u(x) = \begin{cases} 1 & : 0 < x < 1 \\ 0 & : \text{else} \end{cases}$$

uniform



We would like to produce y from x by:

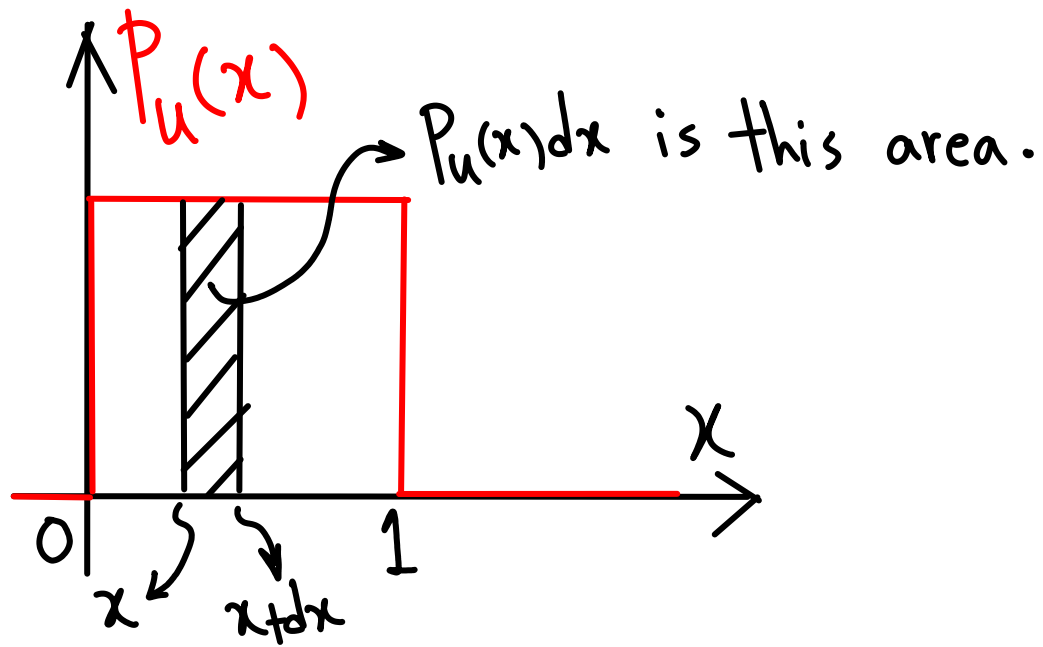
$$y = f(x)$$

in such a way that y possess an arbitrary distribution function:

$$g(y).$$

g can be Gaussian or

How?



Probability of choosing a random number between $(x, x+dx)$: $P_u(x) dx$
 Corresponds to the:
 probability of having a random number y between $(y, y+dy)$: $g(y) dy$

$$P_u(x) dx = g(y) dy$$

$$\int_{-\infty}^x P_u(x) dx = \int_{-\infty}^{y=f(x)} g(y) dy$$

We assume that we know how to perform the integration.

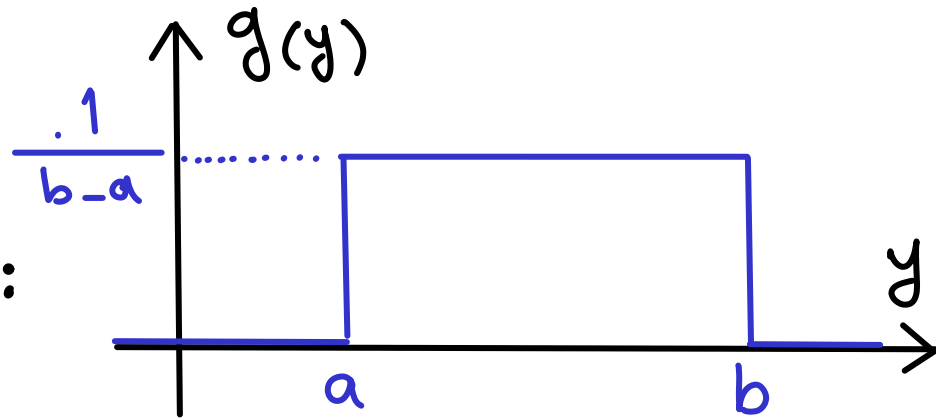
$$x = G(y)$$



$$y = G^{-1}(x)$$

Examples:

- Suppose $g(y)$ has this form:



$$g(y) = \begin{cases} \frac{1}{b-a} & a < y < b \\ 0 & \text{else} \end{cases}$$

Indeed, we choose a random number in $(0, 1)$ and transform it to a random number in (a, b) . How can we do this?

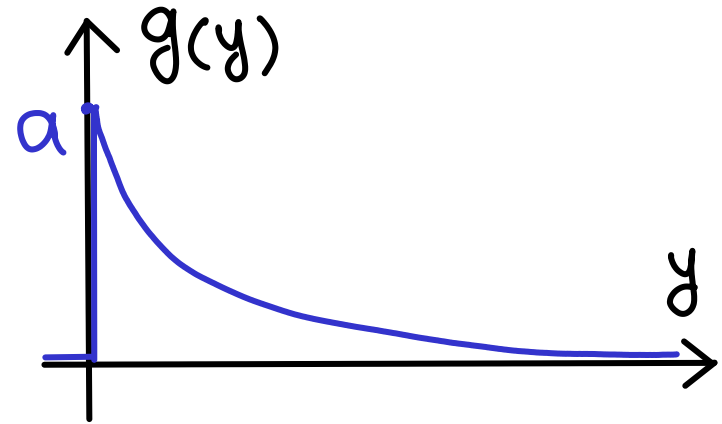
answer:

$$x = \int_{-\infty}^y g(y) dy \Rightarrow x = \int_a^y \frac{1}{b-a} dy \Rightarrow x = \frac{1}{b-a} y \Big|_a^y \Rightarrow x = \frac{y-a}{b-a}$$

$$y = a + (b-a)x$$

Examples:

$$\bullet g(y) = \begin{cases} a e^{-ay} & y > 0 \\ 0 & \text{else} \end{cases}$$



answer:

$$x = \int_{-\infty}^y g(y) dy \Rightarrow x = \int_0^y a e^{-ay} dy \Rightarrow x = a \left(-\frac{1}{a}\right) e^{-ay} \Big|_0^y \Rightarrow x = 1 - e^{-ay}$$

$$e^{-ay} = 1 - x \Rightarrow y = \frac{-1}{a} \ln(1-x) \Rightarrow y = \frac{1}{a} \ln(1-x)^{-1}$$

$$\int_{-\infty}^x P_u(x) dx = \int_{-\infty}^{y=f(x)} g(y) dy$$

We assume that we know how to perform the integration.

$$x = G(y)$$

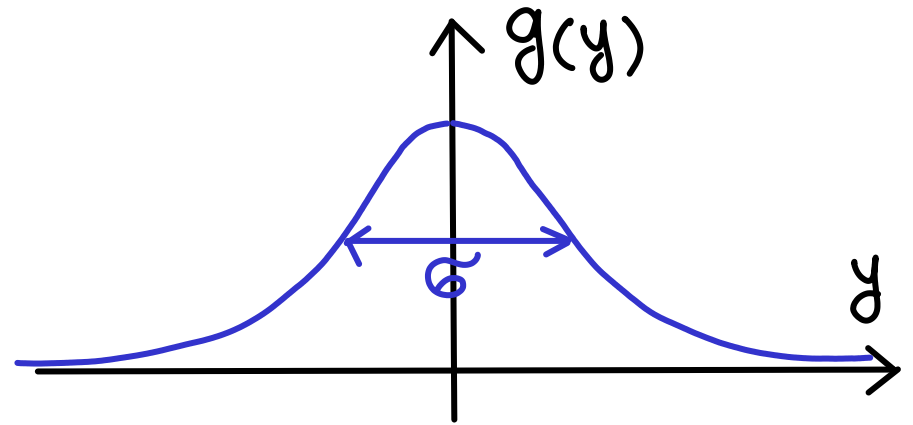


$$y = G^{-1}(x)$$

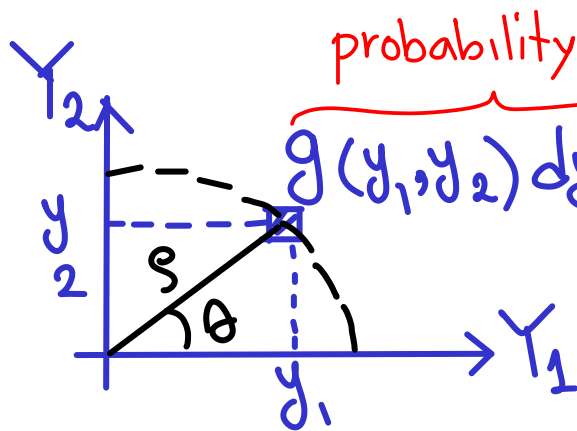
As far as g is integratable and G has inverse this method works.

Examples:

- $g(y) = \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{y^2}{2\sigma^2}}$



answer: $\int_{-\infty}^{\infty} g(y) dy = ?$



probability of having (y_1, y_2)

$$g(y_1, y_2) dy_1 dy_2 = g(y_1) g(y_2) dy_1 dy_2 = \frac{1}{2\pi\sigma^2} e^{-\frac{(y_1^2 + y_2^2)}{2\sigma^2}} dy_1 dy_2$$

if they are independent

$$\sim \frac{1}{2\pi\sigma^2} e^{-r^2/(2\sigma^2)} r dr d\theta = g_\theta(\theta) g_r(r) dr d\theta$$

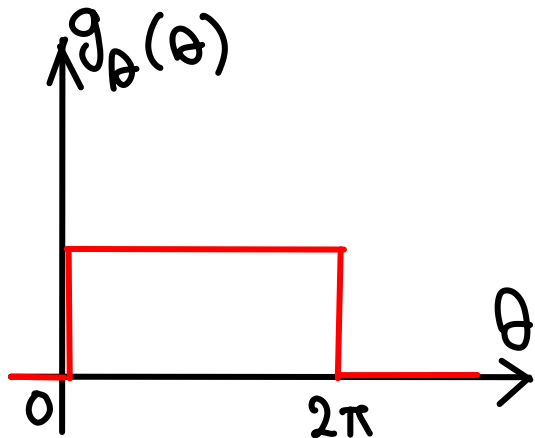
$$\begin{cases} g_\theta(\theta) = \frac{1}{2\pi} \\ g_r(r) = \frac{1}{\sigma^2} r e^{-r^2/(2\sigma^2)} \end{cases}$$

$$P_u(x_1) P_u(x_2) dx_1 dx_2 = g(y_1) g(y_2) dy_1 dy_2$$

$$\int_{-\infty}^{x_1} P_u(x_1) dx_1 \int_{-\infty}^{x_2} P_u(x_2) dx_2 = \int_{-\infty}^{y_1} g(y_1) dy_1 \int_{-\infty}^{y_2} g(y_2) dy_2$$

x_1 (circled in red) and x_2 (circled in blue) are mapped to the following integrals:

 $\int_{\theta=0}^{2\pi} g_{\theta}(\theta) d\theta$ (red box) and $\int_{s=-\infty}^{\infty} \frac{1}{\sigma^2} s e^{-s^2/(2\sigma^2)} ds$ (blue box)



$$a = 0$$

$$b = 2\pi$$

$$\Rightarrow \underbrace{y}_{\theta} = a + (b-a)x_1 \Rightarrow$$

$$\theta = 2\pi x_1$$

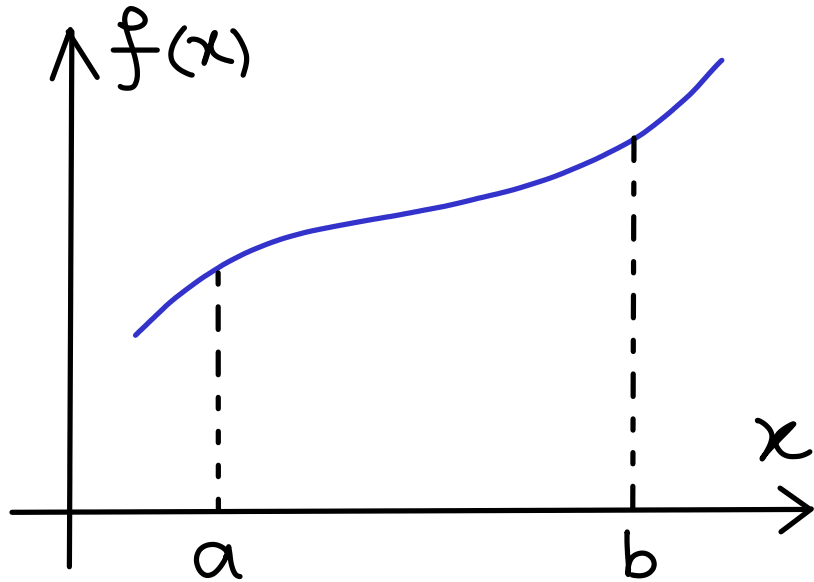
$$\chi_2 = \int_0^{\infty} \frac{1}{\sigma^2} s e^{-\underbrace{s^2/(2\sigma^2)}_y} ds \Rightarrow \chi_2 = \int_{z=s^2/(2\sigma^2)}^{\infty} dz e^{-z} = -e^{-z} \Big|_{z=\frac{s^2}{2\sigma^2}}^{z=\infty}$$

$$\chi_2 = e^{-s^2/(2\sigma^2)} \Rightarrow \boxed{S = +2\sigma^2 \ln \chi_2^{-1}}$$

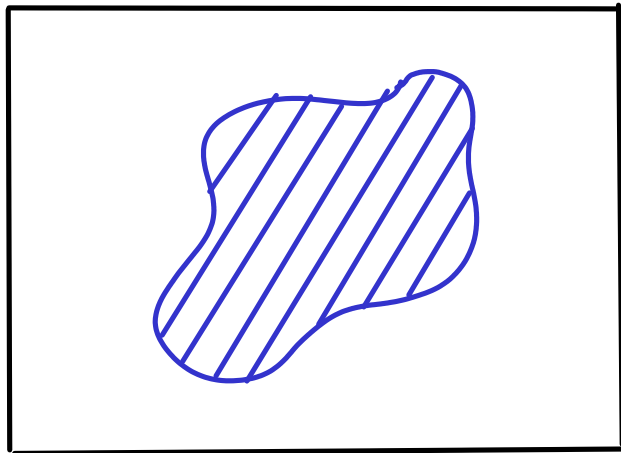
$$\begin{cases} y_1 = S \cos \theta \\ y_2 = S \sin \theta \end{cases}, \begin{cases} \theta = 2\pi \chi_1 \\ S = +2\sigma^2 \ln \chi_2^{-1} \end{cases} \Rightarrow \boxed{\begin{aligned} y_1 &= +2\sigma^2 \cos(2\pi \chi_1) \ln \chi_2^{-1} \\ y_2 &= +2\sigma^2 \sin(2\pi \chi_1) \ln \chi_2^{-1} \end{aligned}}$$

Numerical Integration: a Stochastic Method

Performing integrals

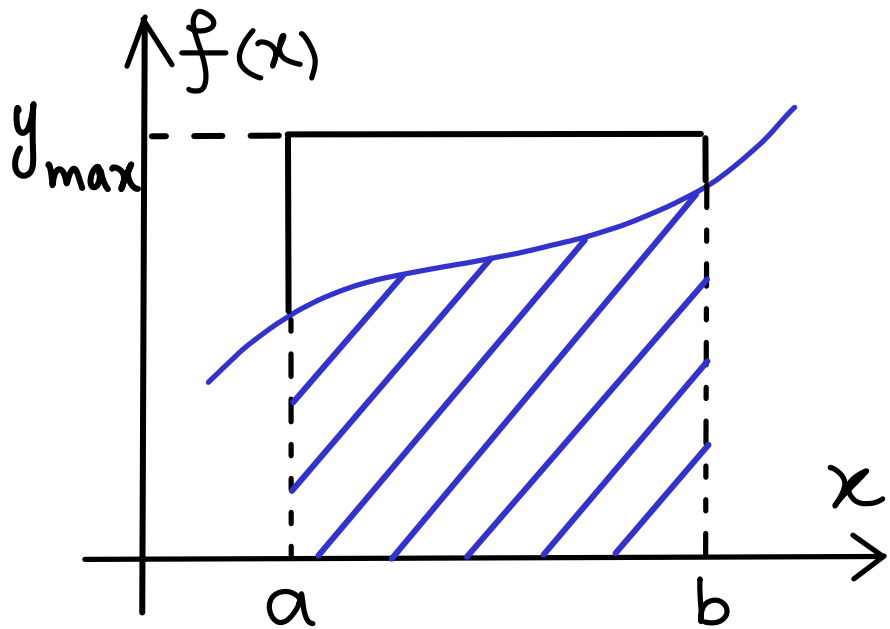


$$I = \int_{x=a}^{x=b} f(x) dx = ?$$



What is the area of the pool inside a garden?

$$A_p = \frac{\text{shelep}}{\text{shelep} + \text{telep}} A_G$$



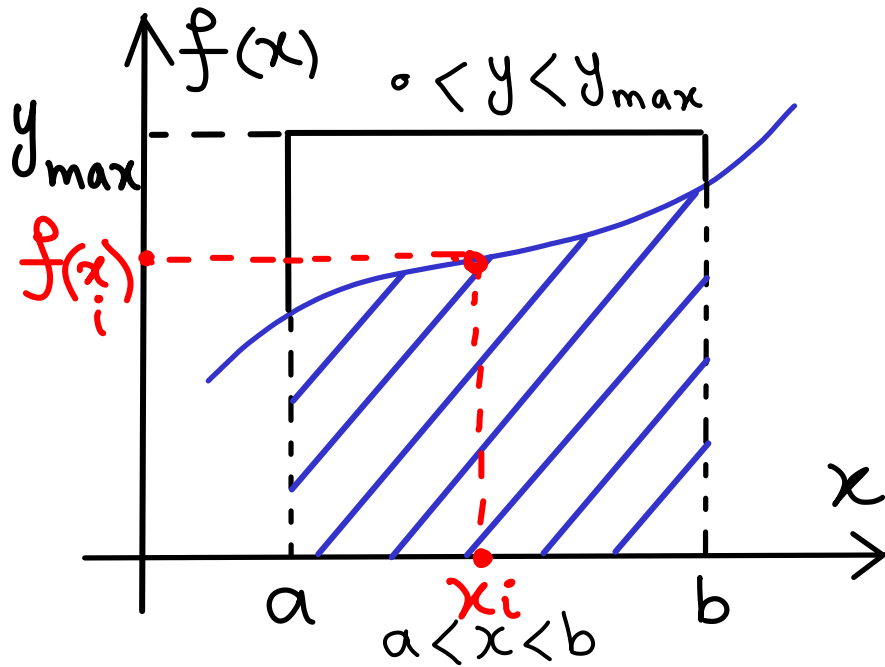
$$a < x < b \quad 0 < y < y_{\max}$$

$$I = \int_{x=a}^{x=b} f(x) dx = ?$$

Loop N | $a < \overset{\text{random}}{x} < b$
 | $0 < y < y_{\max}$
 | If $\overset{\text{random}}{f(x)} > y$ then shelep += 1;

$$\Rightarrow I = \overbrace{(b-a)}^{AG} y_{\max} \frac{\text{Shelep}}{N}$$

This is not a fast algorithm.



$$P(x) = \begin{cases} 1 & a < x < b \\ 0 & \text{else} \end{cases}$$

↓
distribution function

$$\int_{-\infty}^{+\infty} P(x) dx = b - a$$

$$I = \int_{x=a}^{x=b} f(x) dx = (b-a) \frac{\int_{-\infty}^{+\infty} f(x) P(x) dx}{\int_{-\infty}^{+\infty} P(x) dx} = (b-a) \langle f \rangle$$

$$\langle f \rangle = \frac{1}{N} \sum_{i=1}^N f(x_i)$$

$$\begin{array}{l} \text{Loop} \\ N \end{array} \left| \begin{array}{l} I = f(\text{Random}(a, b)) + I; \\ I = \frac{I}{N} * (b-a) \end{array} \right.$$

Simple sampling
Monte Carlo

fluctuations of f

$$\Delta_{\text{error}} = \frac{\sigma_f}{\sqrt{N}}, \quad \sigma_f = \sqrt{\langle f^2 \rangle - \langle f \rangle^2}$$

number of samples

During the run of the code we can find Δ and control the error.

$$I = \int_{x=a}^{x=b} f(x) dx = (b-a) \frac{\int_{-\infty}^{+\infty} f(x) p(x) dx}{\int_{-\infty}^{+\infty} p(x) dx} = (b-a) \langle f \rangle$$

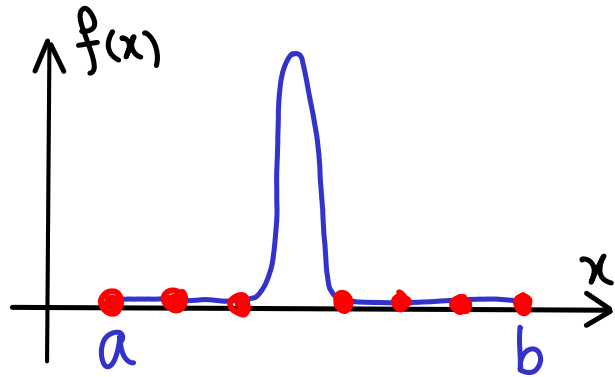
$$\Delta = \frac{\sigma_f}{\sqrt{N}}$$

How error (Δ) can be reduced?

- by increasing N .
- by reducing σ_f . \longrightarrow How?

$$\sigma_f = \sqrt{\langle f^2 \rangle - \langle f \rangle^2}$$

We must reduce $\langle f^2 \rangle - \langle f \rangle^2$.



If the number of sampling is not enough we may miss the peak, that contains the main contribution to the integration.

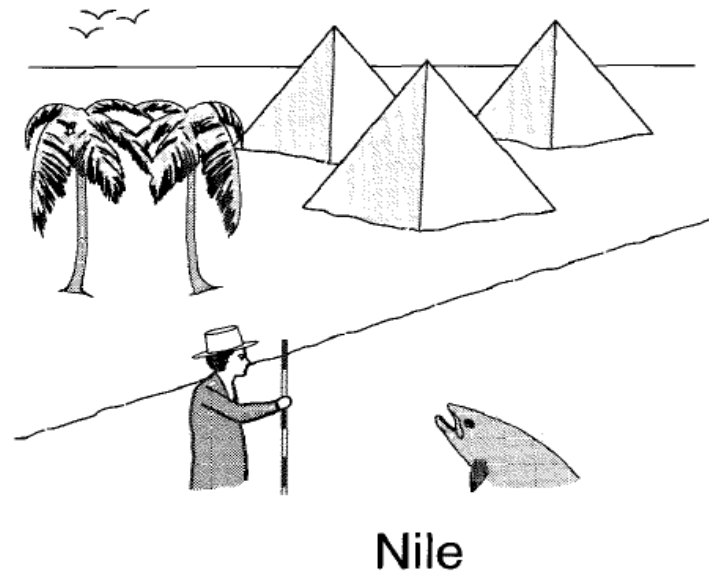
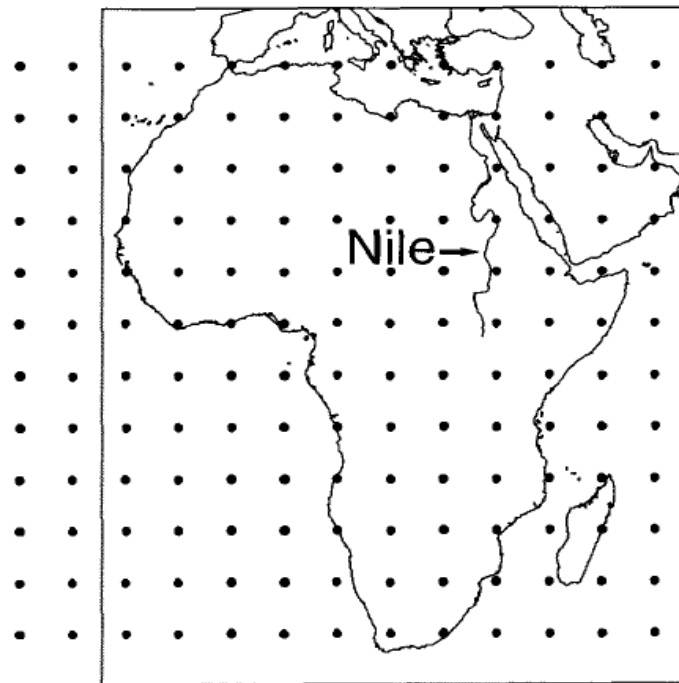


Figure 3.1: Measuring the depth of the Nile: a comparison of conventional quadrature (left), with the Metropolis scheme (right).

Frenkel and Smit, page 28.

$$I = \int_{x=a}^{x=b} f(x) dx = (b-a) \frac{\int_{-\infty}^{+\infty} f(x) p(x) dx}{\int_{-\infty}^{+\infty} p(x) dx} = (b-a) \langle f \rangle$$

$$\Delta = \frac{\tilde{\sigma}_f}{\sqrt{N}}$$

$$\tilde{\sigma}_f = \sqrt{\langle f^2 \rangle - \langle f \rangle^2}$$

To reduce $\tilde{\sigma}_f$, i.e. reducing $\langle f^2 \rangle - \langle f \rangle^2$, we must pick a non-uniform $p(x)$.

$$I = \int_{x=a}^{x=b} f(x) dx = \int g(x) dx \frac{\int \frac{f(x)}{g(x)} g(x) dx}{\int g(x) dx}$$

$$I = \int_{x=a}^{x=b} f(x) dx = \int g(x) dx \frac{\int \frac{f(x)}{g(x)} g(x) dx}{\int g(x) dx}$$

$$I = \int_{x=a}^{x=b} f(x) dx = \int g(x) dx \left\langle \frac{f}{g} \right\rangle_g$$

We know how to perform

averaging over distribution function g

g is not uniform. Therefore, when $g(x)$ is big the probability of choosing x must be big and vice versa.

$$I = \int_{x=a}^{x=b} f(x) dx = \int g(x) dx \left\langle \frac{f}{g} \right\rangle_g$$

$$\Delta = \int g(x) dx \frac{\sigma_{f/g}}{\sqrt{N}} \quad \sigma_{f/g} = \sqrt{\left\langle \left(\frac{f}{g} \right)^2 \right\rangle - \left\langle \frac{f}{g} \right\rangle^2}$$

Now the error Δ depends on the fluctuations of f/g .

This tells us that g should be chosen in such a way that f/g has small fluctuations. This happens if g behaves similar to f .

What is the meaning of $\langle \frac{f}{g} \rangle_g$ in algorithmic language?

Loop
N

$$\left. \begin{array}{l} x = \text{Random} \leftarrow \text{uniform} \\ y = \text{Random}(g(x)) \\ I = f(y)/g(y) + I \end{array} \right|$$

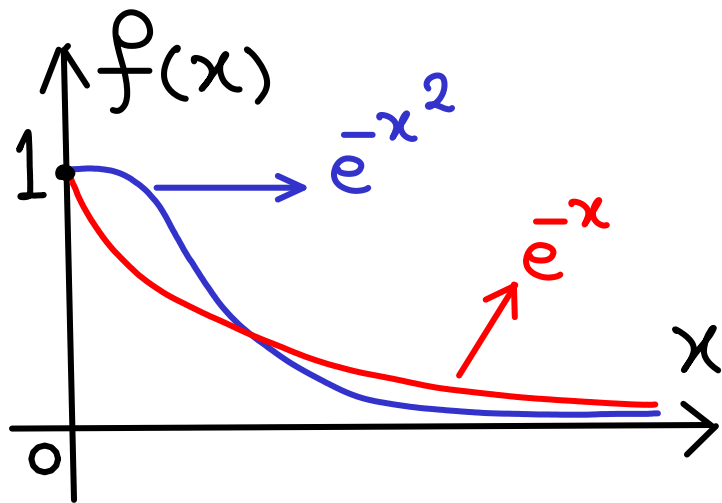
$$I = \frac{I}{N} * \int g(x) dx \leftarrow \text{We knew how to perform this.}$$

Importance Sampling Monte Carlo

Exercise :

Calculate the following integral with both simple sampling and importance sampling Monte Carlo. For the importance sampling one can use g as $g(x) = e^{-x}$.

$$I = \int_0^2 e^{-x^2} dx$$



N	SS				IS	
	I	ΔI	ΔI_m	t		
100						
200						
500						
1000						